Exponentiation Ciphers

Exponentiation ciphers are more resistant to cryptoanalysis.

Let p be an odd prime and let $r \in \mathbb{N}$ with (r, p-1) = 1 be the enciphering key. To encipher,

(1) transform letters to two digit numbers:

- (2) group the resulting numbers into blocks of 2m digits, where m is the largest natural number such that any 2m digit number which could appear is less than p (for example, if 2525
- (3) for each plaintext block x, an integer with 2m digits, form a ciphertext block y by taking y ∈ N with 0 ≤ y

$$y \equiv x^r \pmod{p}$$
.

To decipher, we need the deciphering key r', which is the inverse of r modulo p-1, i.e.

$$rr' \equiv 1 \pmod{p-1}$$
.

Indeed, it follows that there exists $k \in \mathbb{Z}$ such that rr' = 1 + k(p-1). Hence, by Fermat's Theorem 14.8,

$$y \equiv x^r \pmod{p}$$

 $\Rightarrow y^{r'} \equiv x^{rr'} \pmod{p} \equiv x^{1+k(p-1)} \pmod{p} \equiv x \cdot (x^{p-1})^k \pmod{p} \equiv x \pmod{p}$.

So $x \equiv y^{r'} \pmod{p}$.

Examples 16.6. (1) Consider p=29 and r=3. Then (r, p-1)=(3, 28)=1, m=1 and r'=19. Hence we have that

For example, for x = 11

$$11^2 = 121 \equiv 5 \pmod{29} \implies x^r = 11^3 \equiv 55 \pmod{29} \equiv 26 \pmod{29} \implies y = 26.$$