

# Coding Theory

## Sheet 6 Solutions

Spring 2014

1. For  $\text{Ham}(r, q)$ ,

$$n = \frac{q^r - 1}{q - 1}, \quad k = n - r, \quad d = 3.$$

A code  $C$  is MDS if  $d = n - k + 1$ . Hence  $\text{Ham}(r, q)$  is MDS when  $3 = r + 1$ ; that is,  $r = 2$ . So, only  $\text{Ham}(2, q)$  is MDS.

2. (a)  $\text{Ham}(2, 3)$  is a  $[4, 2]$  code.

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 2 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & 2 & 0 & 1 \end{bmatrix} = H.$$

Hence

$$G = \begin{bmatrix} 1 & 0 & 2 & 2 \\ 0 & 1 & 2 & 1 \end{bmatrix}.$$

(b)  $\text{Ham}(2, 4)$  is a  $[5, 3]$  code;  $\mathbf{F}_4 = \{0, 1, \omega, \bar{\omega} \mid \bar{\omega} = \omega + 1 = \omega^2\}$ .

$$\begin{bmatrix} 0 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & \omega & \bar{\omega} \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 1 & 1 & 0 \\ 1 & \omega & \bar{\omega} & 0 & 1 \end{bmatrix} = H.$$

Hence

$$G = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & \omega \\ 0 & 0 & 1 & 1 & \bar{\omega} \end{bmatrix}.$$

(c)  $\text{Ham}(3, 3)$  is a  $[13, 10]$  code. Similarly, to the previous examples, let

$$H = \begin{bmatrix} 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 1 & 1 & 2 & 2 & 2 & 0 & 1 & 0 \\ 1 & 2 & 1 & 2 & 0 & 1 & 2 & 0 & 1 & 2 & 0 & 0 & 1 \end{bmatrix}.$$

Hence

$$H = \left[ \begin{array}{cccccccccc|ccc} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & 2 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & 0 & 2 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & 2 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & 2 & 2 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 2 & 2 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 2 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 2 & 1 & 2 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 2 & 1 & 1 \end{array} \right].$$

(d) Ham(3, 4) is a [21, 18] code. Here,

$$H = \left[ \begin{array}{ccccccccccccccccccccccc} 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & \omega & \omega & \omega & \omega & \bar{\omega} & \bar{\omega} & \bar{\omega} & \bar{\omega} \\ 1 & 0 & 1 & \omega & \bar{\omega} & 0 & 1 & \omega & \bar{\omega} & 0 & 1 & \omega & \bar{\omega} & 0 & 1 & \omega & \bar{\omega} & 0 & 1 & \omega & \bar{\omega} \end{array} \right],$$

$$H' = \left[ \begin{array}{cccccccccccccccccccccccc} 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & \omega & \omega & \omega & \omega & \bar{\omega} & \bar{\omega} & \bar{\omega} & \bar{\omega} & 0 & 1 & 0 \\ 1 & \omega & \bar{\omega} & 1 & \omega & \bar{\omega} & 0 & 1 & \omega & \bar{\omega} & 0 & 1 & \omega & \bar{\omega} & 0 & 1 & \omega & \bar{\omega} & 0 & 0 & 1 \end{array} \right]$$

$$= [B \ I_3].$$

Then  $G' = [I_{18} \ B^T]$ .

(e) Ham(3, 5) is a [31, 28] code. Here,  $H = [A_1 \ A_2] = [B \ I_3]$ , where

$$A_1 = \left[ \begin{array}{cccccccccccccccc} 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 3 & 4 & 1 & 2 & 3 & 4 & 0 & 1 & 2 & 3 & 4 & \end{array} \right]$$

$$A_2 = \left[ \begin{array}{ccccccccccccccccccccccc} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 2 & 2 & 3 & 3 & 3 & 3 & 3 & 4 & 4 & 4 & 4 & 4 & 0 & 1 & 0 \\ 0 & 1 & 2 & 3 & 4 & 0 & 1 & 2 & 3 & 4 & 0 & 1 & 2 & 3 & 4 & 0 & 0 & 1 \end{array} \right].$$

So  $G = [I_{28} \ -B^T]$ .

(f) Ham(4, 2) is a [15, 11] code.

$$\begin{aligned}
 H &= \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \end{bmatrix} \\
 &\rightarrow \begin{bmatrix} 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 1 \end{bmatrix} \\
 &= H' = [B \ I_4].
 \end{aligned}$$

Then  $G = [I_{11} \ B^T]$ .

3. Use  $H$  as in 3(f). A coset leader is 0 or  $l_i$ ,  $i = 1, \dots, 15$ . For a received message  $y_j$ , let  $x_j$  be the corrected message.

(a)  $y_1 = 00000\ 00000\ 11111 \Rightarrow y_1 H^T = 1011 \Rightarrow$  coset leader is  $l_{11}$   
 $\Rightarrow x_1 = 00000\ 00000\ 01111;$

(b)  $y_2 = 00000\ 11111\ 11111 \Rightarrow y_2 H^T = 0001 \Rightarrow$  coset leader is  $l_1$   
 $\Rightarrow x_2 = 10000\ 11111\ 11111;$

(c)  $y_3 = 11111\ 11111\ 11111 \Rightarrow y_3 H^T = 0000 \Rightarrow x_3 = y_3.$

4. Let  $\mathbf{F}_4 = \{0, 1, \omega, \bar{\omega} \mid \bar{\omega} = \omega + 1 = \omega^2\}$ . Use  $H$  as in 3(d).

(a)  $y = 1111111\ 1111111\ 1111111 \Rightarrow y H^T = 001 \Rightarrow x = 0111111\ 1111111\ 1111111;$

(b)  $y' = 1111111\ \omega\omega\omega\omega\omega\omega\ \bar{\omega}\bar{\omega}\bar{\omega}\bar{\omega}\bar{\omega}\bar{\omega} \Rightarrow y' H^T = 1\omega\omega$   
 $\Rightarrow x' = 1111111\ \omega\omega\omega\omega\omega\omega\ \bar{\omega}\bar{\omega}\bar{\omega}\bar{\omega}\bar{\omega}\bar{\omega}.$

5. Since  $x \cap y = (x_1 y_1, \dots, x_n y_n)$ , so

$$x \cap y = (z_1, \dots, z_n), \text{ where } z_i = 1 \Leftrightarrow x_i = y_i = 1.$$

Also

$$x + y = (t_1, \dots, t_n), \text{ with } t_i = 1 \Leftrightarrow (x_i, y_i) = (1, 0) \text{ or } (0, 1).$$

So

$$\begin{aligned}
 w(x + y) &= \text{number of 1's in } x \\
 &\quad + \text{number of 1's in } y \\
 &\quad - 2 \times \text{number of } i \text{ where } x_i = y_i = 1 \\
 &= w(x) + w(y) - 2w(x \cap y).
 \end{aligned}$$

6. Let  $H = [c_1 \cdots c_n]$  in columns; also  $Hx^T = 0$  for  $x \in C$ . So  $Hx^T = x_1c_1 + \cdots + x_nc_n = 0$ . Let  $x' = [x \ x_{n+1}]$ . Then, if  $x \in C$ , we have  $x' \in C'$  when  $x_1 + \cdots + x_n + x_{n+1} = 0$ . So

$$\begin{aligned} H'x'^T &= \begin{bmatrix} c_1 & \cdots & c_n & z^T \\ 1 & \cdots & 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_n \\ x_{n+1} \end{bmatrix} \\ &= \begin{bmatrix} x_1c_1 + \cdots + x_nc_n + 0 & 0 \\ x_1 + \cdots + x_n + x_{n+1} & x_{n+1} \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}. \end{aligned}$$

Hence  $H'$  is a parity-check matrix for  $C'$ .

7. Let  $H = [h_1, \dots, h_r]^T$  be a parity-check matrix of  $C = \text{Ham}(r, 2)$  with rows  $h_1, \dots, h_r$ , and let  $h = \sum \lambda_i h_i$  be an element of  $C^\perp$ . If  $(x_1, \dots, x_r)^T$  is the  $j$ -th column of  $H$ , then the  $j$ -th coordinate of  $h$  is zero if  $\sum \lambda_i x_i = 0$ . However, the number of non-zero solutions  $(x_1, \dots, x_r)$  of  $\sum \lambda_i x_i = 0$  is the number  $N$  of elements of  $V(r, 2) \setminus \{0\}$  in a subspace of dimension  $r - 1$ . Hence  $N = 2^{r-1} - 1$ . So

$$\begin{aligned} w(h) = n - N &= (2^r - 1) - (2^{r-1} - 1) \\ &= 2^r - 2^{r-1} = 2^{r-1}. \end{aligned}$$