

# Coding Theory

## Sheet 5 Solutions

Spring and Summer 2010

1. With

$$G = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 & 1 \end{bmatrix},$$

the following table is calculated.

Symbol	→	$u$	→	$uG$
space	→	0000	→	0000000
A	→	0001	→	1101001
B	→	0010	→	0101010
C	→	0100	→	1001100
D	→	1000	→	1110000
E	→	0011	→	1000011
F	→	0110	→	1100110
G	→	1100	→	0111100
M	→	1001	→	0011001
N	→	1010	→	1011010
O	→	0101	→	0100101
R	→	0111	→	0001111
S	→	1110	→	0010110
T	→	1011	→	0110011
U	→	1101	→	1010101
Y	→	1111	→	1111111

As the minimum weight of a codeword is 3, so  $d(C) = 3$ . Hence  $C$  is single-error correcting. The decoded messages are as follows:

- (a) DECODER
- (b) SEE YOU SOON
- (c) MASTERS STUDENT

2. (a) If  $C = C^\perp$ , then  $n - k = k$ ; so  $k = \frac{1}{2}n$ .  
 (b) Let  $C = \{0000, 1111\}$ . Then  $C$  is a  $[4, 1]_2$  code and  $C^\perp$  is a  $[4, 3]_2$  code;  
 $C^\perp = \{0000, 1111, 0011, 1100, 0101, 1010, 0110, 1001\}$ .
3. (a) As  $|V(5, 2)| = 32$ ,  $|C| = 8$ , there are four cosets.

$C$	00000	10011	01011	00101	11000	10110	01110	11101
$10000 + C$	10000	00011	11011	10101	01000	00110	11110	01101
$00100 + C$	00100	10111	01111	00001	11100	10010	01010	11001
$00010 + C$	00010	10001	01001	00111	11010	10100	01100	11111

- (b) If  $G = [I_3 \ A]$ , then a parity-check matrix is  $H = [A^T \ I_2]$ . Hence

$$H = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 \end{bmatrix}.$$

4. (a) As  $|V(3, 3)| = 27$ ,  $|C| = 9$ , there are three cosets.

$C$	000	210	120	012	021	222	201	102	111
$100 + C$	100	010	220	112	121	022	001	202	211
$200 + C$	200	110	020	212	221	122	101	002	011

- (b) By row operations,

$$G = \begin{bmatrix} 2 & 1 & 0 \\ 0 & 1 & 2 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 2 & 0 \\ 0 & 1 & 2 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 0 & 2 \\ 0 & 1 & 2 \end{bmatrix}.$$

Hence  $H = [1 \ 1 \ 1]$ .